

## Runge-Kutta stepping schemes

### Classical 4-step Runge-Kutta

To solve the equation

$$\frac{dy}{dt} = f(t, y)$$

We use the following scheme

$$k_1 = \Delta t f(t, y)$$

$$k_2 = \Delta t f\left(t + \frac{\Delta t}{2}, y + k_1 / 2\right)$$

$$k_3 = \Delta t f\left(t + \frac{\Delta t}{2}, y + k_2 / 2\right)$$

$$k_4 = \Delta t f(t + \Delta t, y + k_3)$$

$$y(t + \Delta t) = y(t) + (\Delta t / 6)(k_1 + 2k_2 + 2k_3 + k_4)$$

→ 1) Method is order 4

(i.e. error is  $O(\Delta t^5)$ )

2) Does not require initialization

3) Requires 4 function evaluations.

(LMFs require only 1 function evaluation to attain any order)

## Generalized Runge-Kutta schemes

An  $s$ -stage Runge-Kutta scheme has the form

$$k_a = \Delta t f\left(t + c_a \Delta t, y + \sum_{b=1}^s A_{ab} k_b\right),$$

$$a \in \{1, 2, \dots, s\}$$

$$y(t + \Delta t) = y(t) + \sum_{a=1}^s b_a k_a$$

where we introduce the simplifying condition

$$c_a = \sum_{b=1}^s A_{ab}$$

→ Butcher tableau

$c_1$	$A_{11} \quad A_{12} \quad \dots \quad A_{1s}$
$c_2$	$A_{21} \quad A_{22} \quad \dots \quad A_{2s}$
$\vdots$	$\vdots \quad \vdots \quad \ddots \quad \vdots$
$c_s$	$A_{s1} \quad A_{s2} \quad \dots \quad A_{ss}$
	$b_1 \quad b_2 \quad \dots \quad b_s$

A brief way to  
represent a  
Runge-Kutta scheme

→ Explicit / Implicit classification

A Runge-Kutta scheme is

- a) Explicit  $\Leftrightarrow A_{ab} = 0, \forall a, b : a \leq b$
- b) Implicit  $\Leftrightarrow$  not Explicit  
(requires solution of non-linear equations)
- c) Diagonally Implicit (DIRK)  $\Leftrightarrow A_{ab} = 0, \forall a, b : a < b$
- d) Singly Diagonally Implicit (SDIRK)  $\Leftrightarrow$

{ DIRK  
 $A_{11} = A_{22} = \dots = A_{ss}$  (diagonal elements are equal).

→ Typical RKF schemes

① 1-stage RK

$$\begin{array}{c|c} 0 & 0 \\ \hline & 1 \end{array} \quad K_1 = \Delta t f(t, y) \quad \text{Euler rule}$$
$$y_{n+1} = y_n + k_1$$

② 2-stage RG

0	0	0
1	1	0
	1	0

(Euler)

$$k_1 = f(t, y)$$

$$k_2 = f(t + \Delta t, y + \Delta t k_1)$$

$$y = y + \Delta t k_1$$

0	0	0
1	1	0
	0	1

$$y = y + \Delta t k_2$$

(midpoint rule)\*

0	0	0
1	1	0
	1/2	1/2

$$y = y + \Delta t \frac{(k_1 + k_2)}{2}$$

(trapezoidal rule)\*

\* These are EXPLICIT, thus different from  
their LMF cousins

③ 3-stage RK

0	0	0	0
2/3	2/3	0	0
2/3	1/3	1/3	0
	1/4	0	3/4

0	0	0	0
1/2	1/2	0	0
1	-1	2	0
	1/6	2/3	1/6

④ 4-stage Runge-Kutta

0	0	0	0	0
$\frac{1}{2}$	$\frac{1}{2}$	0	0	0
$\frac{1}{2}$	0	$\frac{1}{2}$	0	0
1	0	0	1	0
	$\frac{1}{6}$	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{1}{6}$

(classical Runge-Kutta)

0	0	0	0	0
$\frac{1}{4}$	$\frac{1}{4}$	0	0	0
$\frac{1}{2}$	0	$\frac{1}{2}$	0	0
1	1	-2	2	0
	$\frac{1}{6}$	0	$\frac{2}{3}$	$\frac{1}{6}$

(another 4-stage RG).

## Order Conditions for RG methods

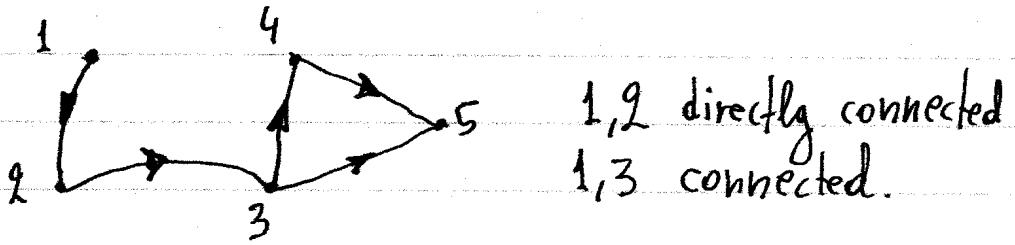
The construction of the order conditions is based on the theory of rooted trees

### → Review of rooted trees

Def : A directed graph  $G = (V, E)$  consists of

- $V =$  a set of vertices
- $E \subseteq V \times V =$  a set of edges

example



$$V = \{1, 2, 3, 4, 5\}$$

$$E = \{(1,2), (2,3), (3,4), (3,5), (4,5)\}$$

Def : Let  $a, b \in V$  be two vertices

- $a, b$  are directly connected iff  $(a, b) \in E$  (notation  $a \Rightarrow b$ )

- $a, b$  are connected iff

$\exists \{x_1, x_2, \dots, x_n\} \subseteq V$  such that

$$a = x_1 \wedge b = x_n \wedge$$

$$\wedge (\forall k \in [n-1] : (x_k, x_{k+1}) \in E)$$

Def : A graph  $G = (V, E)$  is a rooted tree with root  $r \in V$  iff

1) If a vertex  $a \in V$  is not the root, only one arrow goes to  $a$ .

$$\forall a \in V - \{r\} : ((b_1, a), (b_2, a) \in E \Rightarrow b_1 = b_2)$$

$$\forall a \in V - \{r\} : \exists b \in V : (b, a) \in E \wedge b \neq a$$

2) No arrow goes to the root

$$\nexists (a, b) \in E : b \neq r$$

3) Every vertex  $a \in V - \{r\}$  is connected to the root

$$\forall a \in V - \{r\} : r \rightarrow a$$

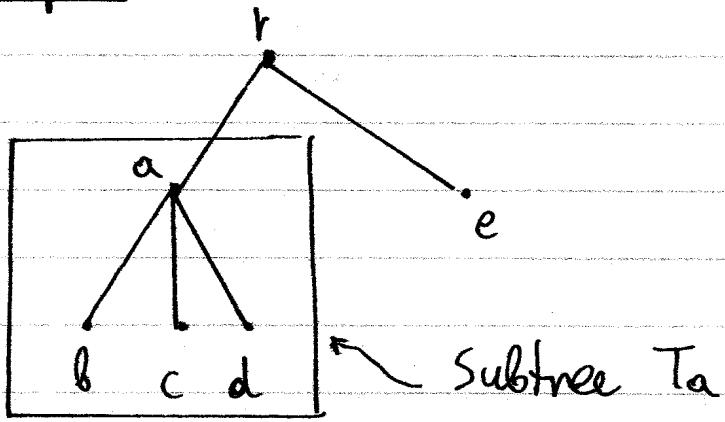
Def : If  $T$  is a rooted tree and  $a \in V - \{r\}$  then  $a$  defines a subtree  $T_a = (V_a, E_a)$

$$V_a = \{x \in V : a \rightarrow x\} \cup \{a\}$$

$$E_a = \{(x, y) \in E : x, y \in V_a\}$$

with root  $a$ .

example



► Tree construction

Let  $\tau$  be the unit tree with  $V = \{r\}$  and  $E = \emptyset$ .

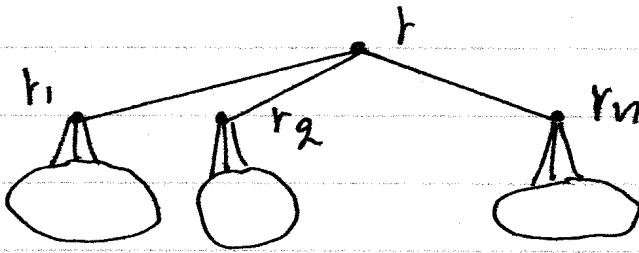
Assume we have constructed the rooted trees

$T_1, T_2, \dots, T_n$ , with roots  $r_1, r_2, \dots, r_n$ .

The tree

$$T = [T_1, T_2, \dots, T_n]$$

is defined as



Any rooted tree can be constructed recursively from  $\tau$  with this operation:

example

$$\{ = [\tau] \quad \vee = [\tau^2]$$

$$\vee = [\tau [\tau^2]] \quad \{ = [[\tau]]$$

► Functions on trees

Let  $T$  be a rooted tree

- 1) The order  $r(T)$  or is  $r(T) = |V| = \text{number of vertices}$
- 2) The density  $\gamma(T)$  is defined as

$$\boxed{\gamma(T) = \prod_{a \in V} r(T_a)}$$

where  $T_a$  is the subtree of  $T$  defined by the vertex  $a$ .

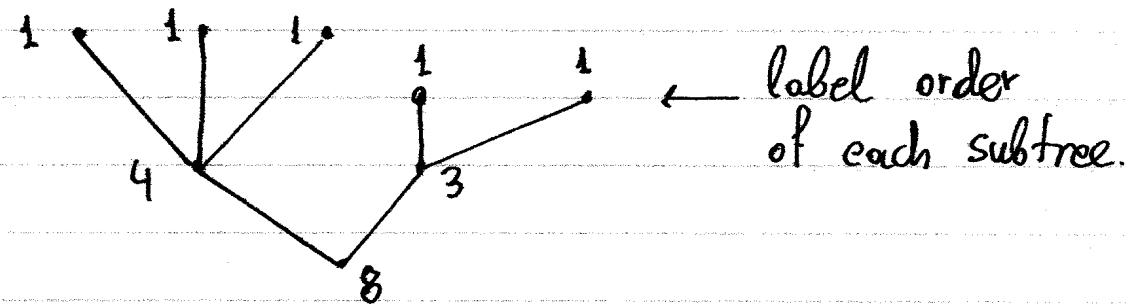
For the unit tree:  $r(\tau) = 1$  and  $\gamma(\tau) = 1$ .

For a tree  $T$  given by

$$T = [T_1^{m_1} T_2^{m_2} \dots T_n^{m_n}]$$

$$\begin{aligned} r(T) &= 1 + \sum_{a=1}^n m_a r(T_a) && \leftarrow \text{order} \\ g(T) &= r(T) \prod_{a=1}^n [g(T_a)]^{m_a} && \leftarrow \text{density.} \end{aligned}$$

example



$$\begin{aligned} g(T) &= 1 \cdot 1 \cdot 1 \cdot 4 \cdot 1 \cdot 1 \cdot 3 \cdot 8 \\ &= 4 \cdot 3 \cdot 8 = 96 \end{aligned}$$

$$r(T) = 8$$



## Order conditions for RG

Consider an RG with tableau

$$\begin{array}{c|c} c_a & A_{ab} \\ \hline & b_a \end{array}, \text{ and } s \text{ steps.}$$

and let  $T$  be a rooted tree.  $\mathbb{T}$  = set of all rooted trees  
We define recursively:

$$(\Phi_a D)(\tau) = 1, \forall a \in [s]$$

$$\Phi_a(T) = \sum_{b=1}^s A_{ab} (\Phi_b D)(T), \forall a \in [s], \forall T \in \mathbb{T}$$

$$(\Phi_a D)([T_1, T_2, \dots, T_m]) = \prod_{b=1}^m \Phi_a(T_b)$$

$$\Phi(T) = \sum_{a=1}^s b_a (\Phi_a D)(T)$$

We call  $\Phi(T)$  = elementary weight of  $T$ .

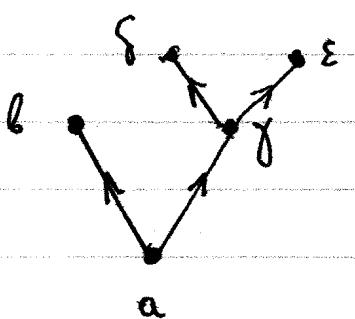
The necessary and sufficient conditions for order  $p$

are

$$1) \sum_{b=1}^s A_{ab} = c_a \quad (\text{simplifying cond.})$$

$$2) \forall T \in \mathbb{T}: (r(T) \leq p \Rightarrow \Phi(T) = \frac{1}{g(T)})$$

## example



The vertices  $b, s, \epsilon$  are called "leaves" bc they have no outgoing arrow.

$$\phi(t) = \sum_{ab\gamma\delta\epsilon} ba A_{ab} A_{\gamma\delta} A_{\gamma\delta} A_{\gamma\epsilon}$$

$$= \sum_{\gamma\delta} ba A_{ab} \left[ \left( \sum_b A_{ab} \right) \left( \sum_\delta A_{\gamma\delta} \right) \left( \sum_\epsilon A_{\gamma\epsilon} \right) \right]$$

$$= \sum_{\gamma\delta} ba A_{ab} c_a c_\gamma c_\delta$$

$$= \sum_{\gamma\delta} ba A_{ab} c_a c_\gamma^2$$

- • 1 The root contributes  $ba$
- 2 Each arrow  $ab$  contributes a factor  $A_{ab}$ .
- 3 The factors contributed by the leaves can be simplified via the simplifying condition.

→ Order conditions for 2-stage RG

$c_1$	0	0
$c_2$	$a_{21}$	0
	$b_1$	$b_2$

$$\left\{ \begin{array}{l} c_1 = 0 \\ c_2 = a_{21} \\ b_1 + b_2 = 1 \\ b_1 c_1 + b_2 c_2 = 1/2 \end{array} \right. \rightarrow \text{simplifying conditions}$$

5 unknowns ; 4 equations ; 1 degree of freedom

→ Order conditions for 3-stage RG

$c_1$	$a_{11}^{10}$	$a_{12}^6$	$a_{13}^6$
$c_2$	$a_{21}$	$a_{22}^0$	$a_{23}^0$
$c_3$	$a_{31}$	$a_{32}^0$	$a_{33}^0$
	$b_1$	$b_2$	$b_3$

$0$	0	0	0
$c_2$	$a_{21}$	0	0
$c_3$	$a_{31}$	$a_{32}$	0
	$b_1$	$b_2$	$b_3$



## Order conditions for 3-stage

0	0	0	0
$c_2$	$a_{21}$	0	0
$c_3$	$a_{31}$	$a_{32}$	0
	$b_1$	$b_2$	$b_3$

$$\left\{ \begin{array}{l} c_2 = a_{21} \\ c_3 = a_{31} + a_{32} \\ b_1 + b_2 + b_3 = 1 \quad [•] \\ b_2 c_2 + b_3 c_3 = 1/2 \quad [!] \\ b_2 c_2^2 + b_3 c_3^2 = 1/3 \quad [√] \\ b_3 a_{32} c_2 = 1/6 \quad [?] \end{array} \right.$$

8 unknowns ; 6 equations.

2 degrees of freedom

Use  $c_2, c_3$  as parameters.

→ Order condition for 4-stage RF

0	0	0	0	0
$c_2$	$a_{21}$	0	0	0
$c_3$	$a_{31}$	$a_{32}$	0	0
$c_4$	$a_{41}$	$a_{42}$	$a_{43}$	0
	$b_1$	$b_2$	$b_3$	$b_4$

$$c_2 = a_{21}$$

$$c_3 = a_{31} + a_{32}$$

$$c_4 = a_{41} + a_{42} + a_{43}$$

$$b_1 + b_2 + b_3 + b_4 = 1 \quad [•]$$

$$b_2 c_2 + b_3 c_3 + b_4 c_4 = 1/2 \quad [?]$$

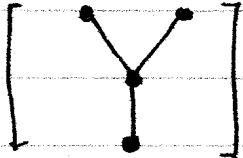
$$b_2 c_2^2 + b_3 c_3^2 + b_4 c_4^2 = 1/3 \quad [V]$$

$$b_3 a_{32} c_2 + b_4 a_{42} c_2 + b_4 a_{43} c_3 = 1/6 \quad [•]$$

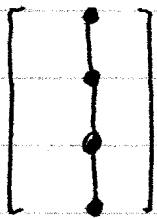
$$b_2 c_2^3 + b_3 c_3^3 + b_4 c_4^3 = 1/4 \quad [V]$$

$$b_3 c_3 a_{32} c_2 + b_4 c_4 a_{42} c_2 + b_4 c_4 a_{43} c_3 = 1/8 \quad [V]$$

$$b_3 a_{32} c_2^2 + b_4 a_{42} c_2^2 + b_4 a_{43} c_3^2 = 1/12$$



$$b_4 a_{43} a_{32} c_2 = 1/94$$



13 unknowns ; 11 equations  
2 degrees of freedom

► It can be shown that  
all possible solutions satisfy  $c_4 = 1$   
One may use  $c_2, c_3$  as parameters.

## Higher-order RG

An explicit RG method of order  $p$  requires the following minimum number of  $s$  stages:

<u>Order p</u>	<u>Minimum s</u>
1, 2, 3, 4	1, 2, 3, 4 (respectively)
5	6
6	7
7	9
8	11
9	12-17 *
10	13-17 *

(\*) Precise minimum unknown.

## Implicit RG

More degrees of freedom in order conditions.  
Current research.  
Excellent stability properties.

## Stability of Runge-kutta schemes

- We have no issues with zero-stability.
- We are interested in absolute stability.

### → The stability polynomial function

Consider the performance of an RG method on the equation

$$\frac{dy}{dt} = qy(t) \text{ with } q \in \mathbb{C}, \operatorname{Re} q < 0$$

We use  $f(t, y) = qy$ . The stages are given by

$$\begin{aligned} k_a &= \Delta t f\left(t + c_a \Delta t, y + \sum_{b=1}^s A_{ab} k_b\right) \\ &= q \Delta t \left(y + \sum_{b=1}^s (A_{ab} k_b)\right) \Leftrightarrow \end{aligned}$$

$$\Leftrightarrow \sum_{b=1}^s (\delta_{ab} - q \Delta t A_{ab}) k_b = (q \Delta t) y$$

$$\Leftrightarrow (I - q \Delta t A) k = (q \Delta t) y \vec{1}$$

with  $k \in \mathbb{R}^s$  a vector and

$$\vec{1} = (1, 1, 1, \dots, 1) \in \mathbb{R}^s$$

This gives

$$K = (I - q\Delta t A)^{-1} q\Delta t y \vec{1}$$

and

$$\begin{aligned}y(t+\Delta t) &\approx y(t) + \sum_{a=1}^s b_a k_a \\&= y + b^T K = \\&= y + b^T (I - q\Delta t A)^{-1} q\Delta t y \vec{1} \\&= [1 + b^T (I - q\Delta t A)^{-1} (q\Delta t y \vec{1})] y\end{aligned}$$

Let  $z = q\Delta t$ . Then  $y(t+\Delta t) = R(z)y(t)$

with

$$R(z) = 1 + b^T (I - zA)^{-1} (z \vec{1})$$

→ stability function.

## Stability Region

Def : The stability region  $\mathcal{S}$  of an RG scheme is defined as

$$\mathcal{S} = \{ z \in \mathbb{C} \mid |R(z)| \leq 1 \}$$

Def : An RG scheme is

- A-stable  $\Leftrightarrow \{z \in \mathbb{C} \mid \operatorname{Re}(z) \leq 0\} \subseteq \mathcal{S}$
- A( $\alpha$ )-stable  $\Leftrightarrow \{z \in \mathbb{C} \mid |\operatorname{Arg}(z) - \pi| < \alpha\} \subseteq \mathcal{S}$
- A(0)-stable  $\Leftrightarrow \exists \alpha \in (0, \pi/2) : A(\alpha)\text{-stable}$

Def : An RG scheme is

$$L\text{-stable} \Leftrightarrow A\text{-stable} \wedge \lim_{\substack{z \rightarrow \infty \\ \operatorname{Re}(z) < 0}} R(z) = 0$$

 An L-stable scheme is even more robust for a very stiff problem than an A-stable scheme.

 For a nonlinear problem, as before, we need  $a(t) \in \mathcal{S}$  for all the eigenvalues  $a$  of the Jacobian  $J(t)$ .

→ Stability function for explicit RG

For an  $s$ -stage explicit RG with order  $p \leq s$   
we note that

- a) The matrix  $A$  is nilpotent with  $A^s = 0$
- b) Rooted trees of the form

$$t_k = \left\{ \begin{array}{c} \uparrow \\ \vdots \\ \uparrow \end{array} \right\} k \text{ vertices}$$

give  $\gamma(t_k) = k!$  and

$$\phi = \sum b_1 A_{12} A_{23} \cdots A_{k-1,k} = b^T A^{k-1} \vec{1}$$

so from the order  $p$  condition we have

$$b^T A^{k-1} \vec{1} = \frac{1}{k!}, \forall k \in \{1, 2, \dots, p\}$$

We use these results to calculate  $R(z)$  as follows:

$$R(z) = 1 + z b^T (I - z A^{-1})^{-1} \vec{1}$$

$$= 1 + z b^T \left[ \sum_{k=0}^{+\infty} (z A)^k \right]^{-1} \vec{1}$$

$$= 1 + z b^T \left[ \sum_{k=0}^{s-1} (z A)^k \right]^{-1} \vec{1}$$

(use nilpotent  $A^s = 0$ )

$$= 1 + z + \sum_{k=1}^{p-1} [b^T A^k \vec{1}] z^{k+1} + z b^T \left[ \sum_{k=p+1}^s (z A)^k \right]^{-1} \vec{1}$$

$$= 1 + z + \sum_{k=2}^p \frac{z^k}{k!} + \sum_{k=p+1}^s c_k z^k$$

thus:

$$R(z) = \sum_{k=0}^p \frac{z^k}{k!} + \sum_{k=p+1}^s c_k z^k$$

$$\text{with } c_k = b^T A^{k-1} \vec{1}$$

We see that

- a) For  $s \leq 4$ , we have  $p=s$  and thus  
 $R(z)$  is universal (i.e. independent of the  
choice of a specific RG scheme)

Specifically:

$$R(z) = \begin{cases} 1+z & , p=1 \\ 1+z + z^2/2 & , p=2 \\ 1+z + z^2/2 + z^3/6 & , p=3 \\ 1+z + z^2/2 + z^3/6 + z^4/24 & , p=4 \end{cases}$$

for ALL explicit RG schemes with order  $p \leq 4$ .

b) For  $p=5$ , we need  $s=6$  and

$$R(z) = 1+z + \frac{z^2}{2} + \frac{z^3}{6} + \frac{z^4}{24} + \frac{z^5}{120} + Cz^6$$

Here  $C$  is dependent on our choice of a particular RG scheme.

c) The stability regions of explicit RG are bounded blobs which expand with increasing  $p$ . This is opposite behaviour from LMF (both explicit/implicit) where the stability region contracts with increasing order  $p$ .

d) See next page for plot of stability regions.

## → Stability Function for Implicit RG

Thm : The stability function can be rewritten

$$R(z) = \frac{N(z)}{D(z)}$$

with  $N(z)$ ,  $D(z)$  polynomials given by

$$N(z) = \det(I + z(\vec{b}^T - A))$$

$$D(z) = \det(I - zA)$$

Def : The polynomial

$E(y) = D(iy)D(-iy) - N(iy)N(-iy)$   
is the E-polynomial of the RG method

Thm : An RG scheme is A-stable iff.

- a) All the zeroes of  $D(z)$  are on  $\{z \in \mathbb{C} \mid \operatorname{Re}(z) > 0\}$
- b)  $E(y) \geq 0$ ,  $\forall y \in \mathbb{R}$ .

Thm : An RG scheme is L-stable iff

- a) It is A-stable
- b)  $\deg N(z) < \deg D(z)$ .

→ To confirm the condition  
 $E(y) \geq 0, \forall y \in \mathbb{R}$

numerically,

- 1 Find the real roots of  $E(y)$
- 2 Check the sign of  $E(y)$  for one point in each interval defined by the real roots!

## Construction of Implicit RG schemes

### ① Gaussian Implicit RG schemes

These are  $s$ -stage schemes with  
order of accuracy  $p=2s$  (!)  
are A-stable  
with tableau

c	A
b	

constructed as follows

- Choose  $c_1, c_2, \dots, c_s$  to be the zeroes of the Legendre polynomial  $P_s(x)$   
(Recall my notes on Gaussian Quadrature)
- Choose  $b_1, b_2, \dots, b_s$  such that they satisfy the linear system

$$\sum_{\alpha=1}^s b_\alpha c_\alpha^{k-1} = \frac{1}{k}, \quad \forall k \in \{1, 2, \dots, s\}$$

which corresponds to the trees:



•<sub>3</sub> Choose  $A_{ab}$  to satisfy

$$\sum_{b=1}^s A_{ab} c_b = \frac{1}{2} c_a^2, \quad \forall a \in \{1, 2, 3, \dots, s\}$$

It can be proved that the resulting RG scheme will indeed

- a) Be A-stable
- b) Have order of accuracy  $p=2s$ .

examples

$\frac{1}{2} - \frac{\sqrt{3}}{6}$	$\frac{1}{4}$	$\frac{1}{4} - \frac{\sqrt{3}}{6}$
$\frac{1}{2} + \frac{\sqrt{3}}{6}$	$\frac{1}{4} + \frac{\sqrt{3}}{6}$	$\frac{1}{4}$
	$\frac{L}{2}$	$\frac{L}{2}$

$$p=2, \quad s=4$$

$\frac{1}{2} - \frac{\sqrt{15}}{10}$	$\frac{5}{36}$	$\frac{2}{9} - \frac{\sqrt{15}}{15}$	$\frac{5}{36} - \frac{\sqrt{15}}{30}$
$\frac{1}{2}$	$\frac{5}{36} + \frac{\sqrt{15}}{24}$	$\frac{2}{9}$	$\frac{5}{36} - \frac{\sqrt{15}}{24}$
$\frac{1}{2} + \frac{\sqrt{15}}{10}$	$\frac{5}{36} + \frac{\sqrt{15}}{30}$	$\frac{2}{9} + \frac{\sqrt{15}}{15}$	$\frac{5}{36}$
	$\frac{5}{18}$	$\frac{4}{9}$	$\frac{5}{18}$

$$p=3, s=6$$

→ We can sacrifice order of accuracy p to attain L-stable implicit R6 schemes

## ② Radau Implicit R6

- Choose  $c_a$  as the roots of  $P_s(x) + P_{s-1}(x)$  with  $P_h(x)$  the Legendre polynomial

- Choose  $b_a, a_{ab}$  same as above

► The resulting method is L-stable with  
 $p = 2s - 1$

example

0	0	0	$s=2$
$\frac{2}{3}$	$\frac{1}{3}$	$\frac{1}{3}$	$p=3$
	$\frac{1}{4}$	$\frac{3}{14}$	

0	0	0	0
$(6-\sqrt{6})/10$	$(9+\sqrt{6})/75$	$(24+\sqrt{6})/120$	$(168-73\sqrt{6})/600$
$(6+\sqrt{6})/10$	$(9-\sqrt{6})/75$	$(168+73\sqrt{6})/600$	$(24-\sqrt{6})/120$
	$\frac{1}{9}$	$(16+\sqrt{6})/36$	$(16-\sqrt{6})/36$

$s=3, p=5$

### ③ Lobatto Implicit RG

- Choose  $a_0$  zeroes of  $P_s(x) - P_{s-2}(x)$   
where  $P_n(x)$  is the Legendre polynomial
- Choose  $b_a, A_{ab}$  same as before.

<u>example</u> :	0	0	0	0	$s=3$
	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$	0	$p=4$
	1	0	1	0	
		$\frac{1}{6}$	$\frac{2}{3}$	$\frac{1}{6}$	

► The resulting method is L-stable with  
 $p = 2s - 32$

#### ④ Diagonally Implicit RK (DIRK)

We sacrifice more order to obtain faster algorithm.

##### 2-stage DIRK

$$\begin{array}{c|cc} \lambda & \lambda & 0 \\ \hline c_2 & c_2 - \lambda & 1 \\ \hline b_1 & b_2 \end{array} \quad \text{with } \begin{cases} b_1 + b_2 = 1 \\ b_1\lambda + b_2 c_2 = 1/2 \end{cases}$$

$$p = 2 ; \quad A\text{-stable} \Leftrightarrow \lambda \geq 1/4 \\ L\text{-stable} \Leftrightarrow \lambda = 1 \pm (1/2)\sqrt{2}$$

##### 3-stage DIRK

$$\begin{array}{c|ccc} \lambda & \lambda & 0 & 0 \\ \hline (3/2)(\lambda+1) & (1/2)(1-\lambda) & \lambda & 0 \\ \hline 1 & (1/4)(-6\lambda^2 + 16\lambda - 1) & (1/4)(6\lambda^2 - 20\lambda + 5) & 1 \\ \hline & (1/4)(-6\lambda^2 + 16\lambda - 1) & (1/4)(6\lambda^2 - 20\lambda + 5) & 1 \end{array}$$

$$p=3$$

A-stable  $\Leftrightarrow \lambda \in [1/3, \alpha]$

L-stable  $\Leftrightarrow \lambda = \beta$

with  $\alpha \approx 1.068$  the root of

$$9x^3 - 3x^2 + x - 1/92 = 0$$

and  $\beta \approx 0.4358$  the root of

$$x^3 - 3x^2 + (3/2)x - 1/6 = 0$$

The advantage of DIRK is that each stage is decoupled from subsequent stages, and can be done by solving 1 nonlinear equation instead of advancing all stages simultaneously by solving a system of  $s$  equations!

## Comparison Chart

<u>LMF</u>	<u>Runge-Kutta</u>
Efficient ; one evaluation of $f$	Less efficient Multiple stages
Difficult to adjust $\Delta t$	Easy to adjust $\Delta t$
Order Barrier $p \leq s+2$ for convergence	No order barriers. A-stability can be achieved for $p=2s$ and L-stability for $p=2s-1$
Order Barrier $p \leq 2$ for A-stability	
Explicit LMF Stability region contracts when $p \nearrow$	Explicit RK Stability region expands when $p \nearrow$